

# Developing Efficient and Portable Communication Software with ACE and C++

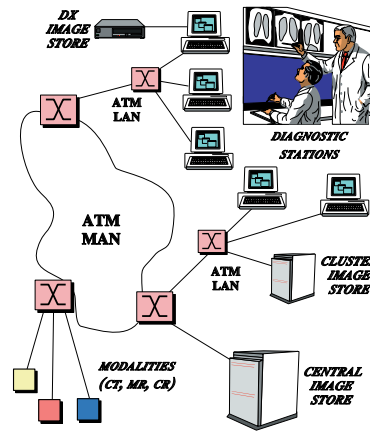
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<http://www.cs.wustl.edu/~schmidt/ACE-examples4.ps.gz>

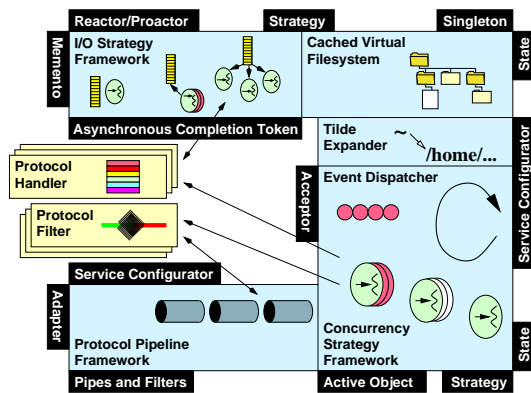
## Problem: Software Evolution



### Key Challenges

- Communication software evolves over time
  - \* Requirements change
  - \* Platforms change
  - \* New design forces emerge
- It is essential to *plan* for inevitable change

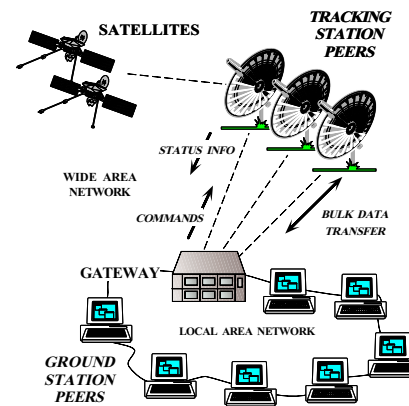
## Solution: Plan for Change Using Frameworks and Patterns



### Solution Approach

- Identify sources of *commonality* and *variability*
- Use patterns to identify reusable design artifacts
- Use frameworks to "unify" variation in code artifacts

## Sources of Variation in Communication Software



### Syntactic Variations

- Unsupported non-essential APIs
- Gratuitous differences in API

### Semantic Variations

- Underlying platform differences
- Framework must respect these differences

### Complex Variations

- Unsupported essential portions of API
- Emulation is necessary

## ACE framework: Resolving Syntactic Variations

```
int ACE_OS::fstat (ACE_HANDLE handle,
                  struct stat *stp)
{
    #if defined (ACE_PSOS_LACKS_PHILE)
        ACE_UNUSED_ARG (handle);
        ACE_UNUSED_ARG (stp);
        ACE_NOTSUP_RETURN (-1);
    #elif defined (ACE_PSOS)
        ACE_OSCALL_RETURN
            (::fstat_f (handle, stp), int, -1);
    #else
        ACE_OSCALL_RETURN
            (::fstat (handle, stp), int, -1);
    #endif /* ACE_PSOS_LACKS_PHILE */
}
```

### • Examples

- *Unsupported*
  - \* Provide “no-op” definitions
  - \* Conditional compilation
- *Syntax*
  - \* Re-map function parameters

## ACE framework: Resolving Semantic Variations

```
int ACE_OS::clock_gettime
    (clockid_t clockid, struct timespec *ts)
{
    #if defined (ACE_HAS_CLOCK_GETTIME)
        ACE_OSCALL_RETURN (::clock_gettime
            (clockid, ts), int, -1);
    #elif defined (ACE_PSOS)
        ACE_UNUSED_ARG (clockid);
        ACE_PSOS_Time_t pt;
        int result = ACE_PSOS_Time_t::get_system_time (pt);
        *ts = ACE_static_cast (struct timespec, pt);
        return result;
    #else
        ACE_UNUSED_ARG (clockid);
        ACE_UNUSED_ARG (ts);
        ACE_NOTSUP_RETURN (-1);
    #endif /* ACE_HAS_CLOCK_GETTIME */
}
```

### • Examples

- *Underlying differences*
  - \* Time in clock ticks
  - \* Ticks-per-second is board-dependent
- *Framework must respect these differences*
  - \* Provide a consistent abstraction
  - \* Intermediate wrappers are useful for small, coherent abstractions

## ACE framework: Resolving Complex Variations

```
void *ACE_TSS_Emulation::tss_open
    (void *ts_storage[ACE_TSS_KEYS_MAX])
{
    #if defined (ACE_PSOS)
        u_long tss_base;
        tss_base = (u_long) ts_storage;
        t_setreg (0, PSOS_TASK_REG_TSS, tss_base);

        void **tss_base_p = ts_storage;
        for (u_int i = 0;
             i < ACE_TSS_KEYS_MAX;
             ++i, ++tss_base_p)
            *tss_base_p = 0;
        return (void *) tss_base;
    #elif defined (...)
        // ...
    #endif
}
```

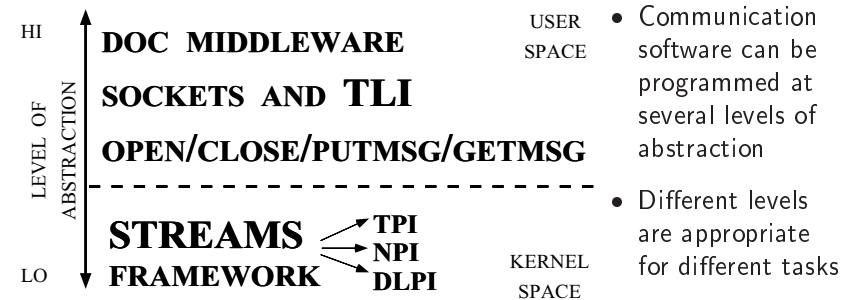
### • Examples

- *Unsupported but essential portions of the API (e.g., thread-specific storage)*
  - \* Provided by POSIX, NT
  - \* Not provided by VxWorks, pSOS

### • Emulation in user space is necessary

- Create a TSS emulation class
- Provide platform-specific method implementations

## Network Programming Alternatives



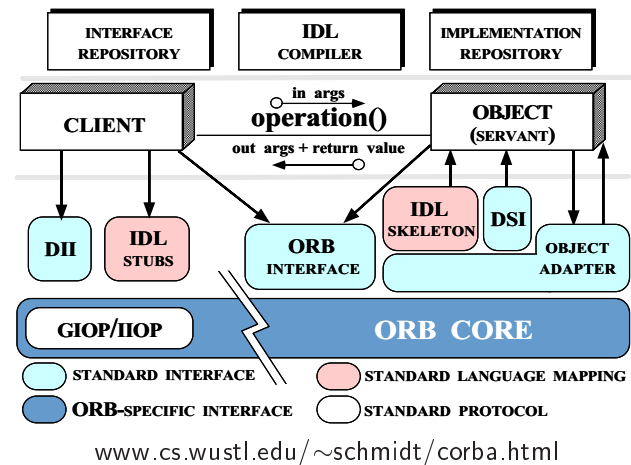
- Communication software can be programmed at several levels of abstraction
- Different levels are appropriate for different tasks

## Navigating Through the Design Alternatives

Choosing the appropriate level of abstraction to program involves many factors

- Performance
  - Higher levels may be less efficient
- Functionality
  - Certain features, e.g., multicast, are not available at all levels
- Ease of programming
  - DOC middleware is typically easier to use
- Portability
  - The socket API is generally portable...

## Overview of DOC Middleware



- Helps simplify many types of applications
- Lets developers work at higher levels of abstraction
- Examples include CORBA, DCOM, Java RMI, DCE, Sun RPC

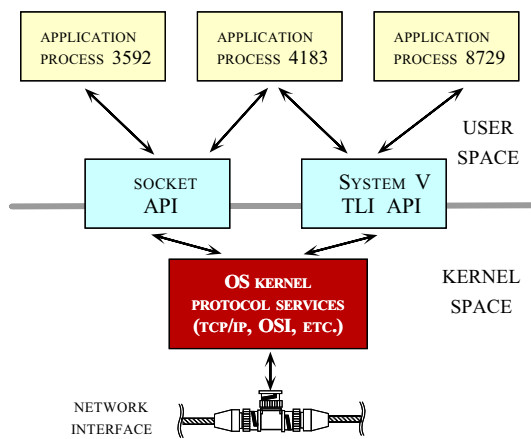
## Common DOC Middleware Features

- DOC middleware “stub/skeleton compiler” support
  - Automatically generate code to perform presentation layer conversions
    - \* e.g., network byte-ordering and parameter marshaling
- DOC middleware runtime support
  - Handle network addressing and remote service identification
  - Perform service registration, port monitoring, and service dispatching
  - Enforce authentication and security
  - Manage transport protocol selection and request delivery
  - Provide reliable operation delivery
  - Demultiplexing and dispatching
  - Concurrency and connection management

## DOC Middleware Limitations

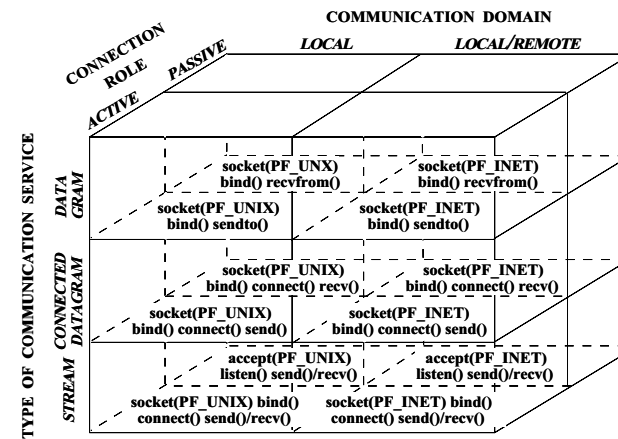
- Some applications may need to access lower-level IPC mechanisms directly to meet certain requirements
  - e.g., performance, functionality, portability, etc.
- Compared with direct use of sockets and TLI, DOC middleware may be less efficient due to
  - Presentation conversion processing and excessive data copying
  - Synchronous client-side and server-side stub behavior
  - Stop-and-wait flow control
  - Non-adaptive retransmission timer schemes
  - Non-optimized demultiplexing and concurrency models

## Standard APIs for Network IPC



- Sockets and TLI allow access to lower-level IPC mechanisms, e.g.:
  - TCP/IP
  - XNS and Novell IPX NetWare protocols
  - UNIX domain sockets
  - OSI protocols

## Socket Taxonomy



- The Socket API can be classified along three dimensions

## Problem with Sockets: Lack of Type-safety

```
int buggy_echo_server (u_short port_num)
{ // Error checking omitted.
  sockaddr_in s_addr;
  int s_fd = socket (PF_UNIX, SOCK_DGRAM, 0);
  s_addr.sin_family = AF_INET;
  s_addr.sin_port = port_num;
  s_addr.sin_addr.s_addr = INADDR_ANY;

  bind (s_fd, (sockaddr *) &s_addr,
        sizeof s_addr);
  int n_fd = accept (s_fd, 0, 0);
  for (;;) {
    char buf[BUFSIZ];
    ssize_t n = read (s_fd, buf, sizeof buf);
    if (n <= 0) break;
    write (n_fd, buf, n);
  }
}
```

- I/O handles are not amenable to strong type checking at compile-time
- The adjacent code contains many subtle, common bugs

## Problem with Sockets: Steep Learning Curve

Many socket/TLI API functions have complex semantics, e.g.:

- Multiple protocol families and address families
  - e.g., TCP, UNIX domain, OSI, XNS, etc.
- Infrequently used features, e.g.:
  - Broadcasting/multicasting
  - Passing open file handles
  - Urgent data delivery and reception
  - Asynch I/O, non-blocking I/O, I/O-based and timer-based event multiplexing

## Problem with Sockets: Poorly Structured

```

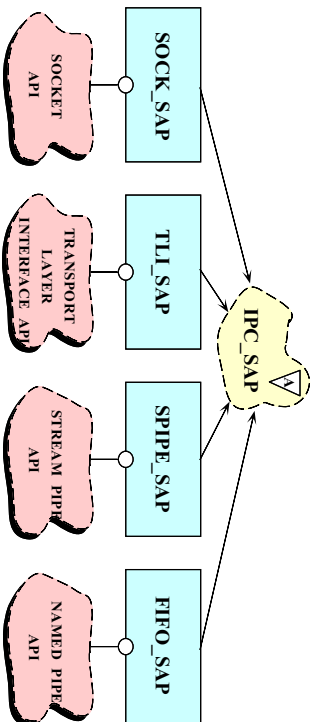
socket()
bind()
connect()
listen()
accept()
read()
write()
readv()
writev()
recv()
send()
recvfrom()
sendto()
recvmsg()
sendmsg()
setsockopt()
getsockopt()
getpeername()
getsockname()
gethostbyname()
getservbyname()

```

- Note the socket API is *linear* rather than *hierarchical*
  - Thus, it gives no hints on how to use it correctly
- In addition, there is no consistency among names...

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## The ACE C++ IPC Wrapper Solution



- ACE provides C++ “wrappers” that encapsulate IPC programming interfaces like sockets and TLI
  - This is an example of the *Wrapper Facade Pattern*

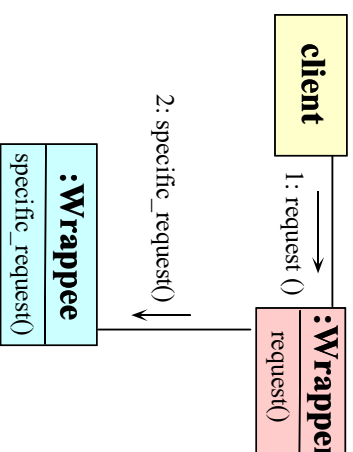
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## Problem with Sockets: Portability

- Having multiple “standards,” *i.e.*, sockets and TLI, makes portability difficult, *e.g.*,
  - May require conditional compilation
  - In addition, related functions are not included in POSIX standards
    - \* *e.g.*, select, WaitForMultipleObjects, and poll
- Portability between UNIX and Win32 Sockets is problematic, *e.g.*:
  - Header files
  - Error numbers
  - Handle vs. descriptor types
  - Shutdown semantics
  - I/O controls and socket options

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## Intent and Structure of the Wrapper Facade Pattern

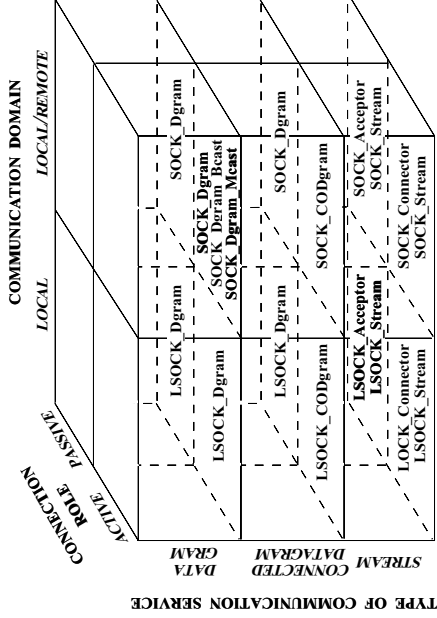


- **Intent**
  - Encapsulates low-level, stand-alone system mechanisms within type-safe, modular, and portable class interfaces
- **Forces Resolved**
  - Avoid tedious, error-prone, and non-portable system APIs
  - Create cohesive abstractions

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## The ACE C++ Socket Wrapper Class Structure

- Note how stand-alone functions are replaced by C++ class components



## SOCK\_SAP Factory Class Interfaces

```
class SOCK_Connector
{
public:
    // Traits
    typedef INET_Addr PEER_ADDR;
    typedef SOCK_Stream PEER_STREAM;

    int connect
        (SOCK_Stream &new_sap,
         const INET_Addr &raddr,
         Time_Value *timeout,
         const INET_Addr &laddr);
    // ...
};

class SOCK_Acceptor
: public SOCK
{
public:
    // Traits
    typedef INET_Addr PEER_ADDR;
    typedef SOCK_Stream PEER_STREAM;

    SOCK_Acceptor
        (const INET_Addr &local_addr);
    int accept
        (SOCK_Stream &new_sap,
         INET_Addr *,
         Time_Value *);
    //...
};
```

## SOCK\_SAP Stream and Addressing Class Interfaces

```
class SOCK_Stream : public SOCK
{
public:
    // Trait.
    typedef INET_Addr PEER_ADDR;

    ssize_t send (const void *buf,
                 int n);
    ssize_t recv (void *buf,
                 int n);
    ssize_t send_n (const void *buf,
                  int n);
    ssize_t recv_n (void *buf,
                  int n);
    int close (void);
    // ...
};

class INET_Addr : public Addr
{
public:
    INET_Addr (u_short port_number,
              const char host[]);
    u_short get_port_number (void);
    int32 get_ip_addr (void);
    // ...
};
```

## OO Design Interlude

Q: *Why decouple the SOCK\_Acceptor and the SOCK\_Connector from SOCK\_Stream?*

A: For the same reasons that Acceptor and Connector are decoupled from Svc\_Handler, e.g.,

- A SOCK\_Stream is only responsible for data transfer
  - Regardless of whether the connection is established passively or actively
- This ensures that the SOCK\* components are not used incorrectly...
  - e.g., you can't accidentally read or write on SOCK\_Connectors or SOCK\_Acceptors, etc.

## ACE C++ Wrapper echo\_server

```
int echo_server (u_short port_num)
{
    // Error handling omitted.
    INET_Addr my_addr (port_num);
    SOCK_Acceptor acceptor (my_addr);
    SOCK_Stream new_stream;

    acceptor.accept (new_stream);

    for (;;)
    {
        char buf[BUFSIZ];
        // Error caught at compile time!
        ssize_t n = acceptor.recv (buf, sizeof buf);
        new_stream.send_n (buf, n);
    }
}
```

## A Generic Version of the Echo Server

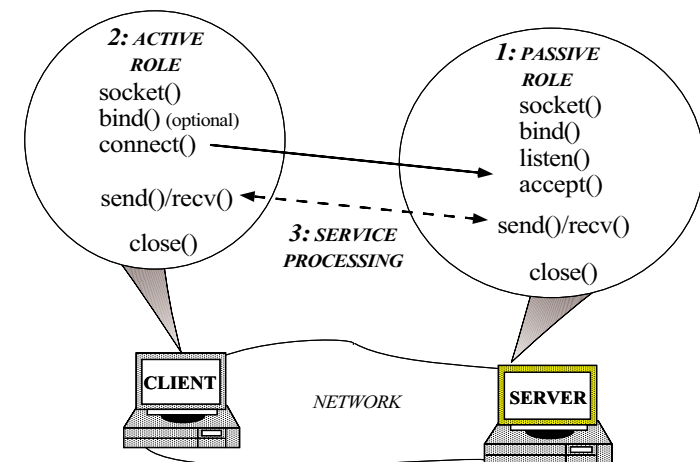
```
template <class ACCEPTOR>
int echo_server (u_short port)
{
    // Local address of server (note use of traits).
    ACCEPTOR::PEER_ADDR my_addr (port);
    // Initialize the passive mode server.
    ACCEPTOR acceptor (my_addr);
    // Data transfer object (note use of traits).
    ACCEPTOR::PEER_STREAM stream;
    // Accept a new connection.
    acceptor.accept (stream);

    for (;;) {
        char buf[BUFSIZ];
        ssize_t n = stream.recv (buf, sizeof buf);
        stream.send_n (buf, n);
    }
}
```

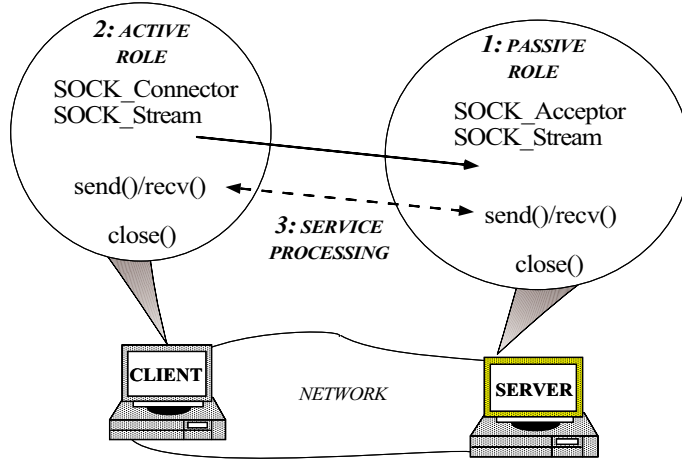
## Socket vs. ACE C++ Socket Wrapper Example

- The following slides illustrate differences between using the Socket interface vs. the ACE C++ Socket wrappers
- The example is a simple client/server “network pipe” application that behaves as follows:
  - Starts an *iterative daemon* at a well-known server port
  - Client connects to the server and transmits its standard input to the server
  - The server prints this data to its standard output
- The server portion of the “network pipe” application may actually run either locally or remotely...

## Network Pipe with Sockets



## Network Pipe with ACE C++ Socket Wrappers



## Running the Network Pipe Program

- e.g.,
 

```
% ./server &
% echo "hello world" | ./client localhost
client localhost.cs.wustl.edu%: hello world
```
- Note that the ACE C++ Socket wrapper example:
  - Requires *much* less code (about 1/2 to 2/3 less)
  - Provides greater clarity and less potential for errors
  - Operates at no loss of efficiency
- Complete example available at URL:
  - [www.cs.wustl.edu/~schmidt/IPC\\_SAP-92.ps.gz](http://www.cs.wustl.edu/~schmidt/IPC_SAP-92.ps.gz)

### Socket Client

```
#define PORT_NUM 10000
int
main (int argc, char *argv[]) {
    struct sockaddr_in saddr;
    struct hostent *hp;
    char *host = argc > 1 ? argv[1] : "tango.cs.wustl.edu";
    u_short port_num = argc > 2
        ? atoi (argv[2]) : PORT_NUM;
    char buf[BUFSIZ];
    int s_fd;
    int w_bytes;
    int r_bytes;
    int n;

    /* Create a local endpoint of communication */
    s_fd = socket (PF_INET, SOCK_STREAM, 0);

    /* Determine IP address of the server */
    hp = gethostbyname (host);
```

### Socket Client (cont'd)

```
/* Set up the address information to
   contact the server */
memset ((void *) &saddr, 0, sizeof saddr);
saddr.sin_family = AF_INET;
saddr.sin_port = port_num;
memcpy (&saddr.sin_addr, hp->h_addr, hp->h_length);

/* Establish connection with remote server */
connect (s_fd, (struct sockaddr *) &saddr,
        sizeof saddr);

/* Send data to server (correctly handles
   "incomplete writes" due to flow control) */
while ((r_bytes = read (0, buf, sizeof buf)) > 0)
    for (w_bytes = 0; w_bytes < r_bytes; w_bytes += n)
        n = write (s_fd, buf + w_bytes, r_bytes - w_bytes);

/* Explicitly close the connection */
close (s_fd);
return 0;
}
```



```

}
}
close (n_fd);
//listening endpoint remains open) */
// Close the new endpoint
write (1, buf, r_bytes);
while ((r_bytes = read (n_fd, buf, sizeof buf)) > 0)
/* Read data from client (terminate on error) */
printf ("client %s: ", hp->h_name), fflush (stdout);
cli_addr_len, AF_INET);
hp = gethostbyname ((char *) &cli_addr.sin_addr,
continue;
if (n_fd == -1)
continue;
&& errno == EINTR)
&cli_addr,
cli_addr,
while ((n_fd = accept (s_fd, (struct sockaddr *)
/* Create a new endpoint of communication */
struct hostent *hp;
int r_bytes, cli_addr_len = sizeof cli_addr;
struct sockaddr_in cli_addr;
char buf[BUFFSZ];
for (;;) {
/* Performs the iterative server activities */

```

## Socket Server (cont'd)

ACE C++ Wrapper Tutorial

```

// Send data to server (correctly handles
// "incomplete writes").
for (;;) {
    ssize_t r_bytes = read (0, buf, sizeof buf);
    cli_stream.send_n (buf, r_bytes);
}
// Explicitly close the connection.
cli_stream.close ();
return 0;
}

```

## C++ Socket Wrapper Client (cont'd)

ACE C++ Wrapper Tutorial

```

#define PORT_NUM 10000
int
main (int argc, char *argv[])
{
    n_short port_num =
        htons (argc < 1 ? atoi (argv[1]) : PORT_NUM);
    struct sockaddr_in saddr;
    int s_fd, n_fd;
    /* Create a local endpoint of communication */
    s_fd = socket (PF_INET, SOCK_STREAM, 0);
    /* Set up the address information to
    become a server */
    memset ((void *) &saddr, 0, sizeof saddr);
    saddr.sin_family = AF_INET;
    saddr.sin_port = port_num;
    saddr.sin_addr.s_addr = INADDR_ANY;
    /* Associate address with endpoint */
    bind (s_fd, (struct sockaddr *) &saddr,
        sizeof saddr);
    /* Make endpoint listen for service requests */
    listen (s_fd, 5);

```

## Socket Server

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```

const n_short PORT_NUM = 10000;
int main (int argc, char *argv[])
{
    char buf[BUFFSZ];
    char *host = argv[1] ? argv[1] : "ics.uct.edu";
    n_short port_num =
        htons (argc < 2 ? atoi (argv[2]) : PORT_NUM);
    INET_Addr server_addr (port_num, host);
    SOCK_Stream cli_stream;
    SOCK_Connector connector;
    // Establish the connection with server.
    connector.connect (cli_stream, server_addr);

```

## C++ Socket Wrapper Client

ACE C++ Wrapper Tutorial

## C++ Wrapper Socket Server

ACE C++ Wrapper Tutorial

```
const u_short PORT_NUM = 10000;
// SOCK_SAP Server.
int
main (int argc, char *argv[])
{
    u_short port_num =
        argc == 1 ? PORT_NUM : atoi (argv[1]);
    // Create a server.
    SOCK_Acceptor acceptor ((INET_Addr) port_num);
    SOCK_Stream new_stream;
    INET_Addr cli_addr;
```

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## C++ Wrapper Socket Server (cont'd)

ACE C++ Wrapper Tutorial

```
// Performs the iterative server activities.
for (;;) {
    char buf[BUFFSZ];
    // Create a new SOCK_Stream endpoint (note
    // automatic restart if errno == EINTR).
    acceptor.accept (new_stream, &cli_addr);
    printf ("Client %s: ", cli_addr.get_host_name ());
    fflush (stdout);
    // Read data from client (terminate on error).
    for (;;) {
        ssize_t r_bytes;
        r_bytes = new_stream.recv (buf, sizeof buf);
        write (1, buf, r_bytes);
    }
    // Close new endpoint (listening
    // endpoint stays open).
    new_stream.close ();
}
```

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ACE C++ Wrapper Tutorial

## ACE C++ Wrapper Design Principles

- Enforce typesafety at compile-time
- Allow controlled violations of typesafety
- Simplify for the common case
- Replace one-dimensional interfaces with hierarchical class categories
- Enhance portability with parameterized types
- Inline performance critical methods
- Define auxiliary classes to hide error-prone details

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ACE C++ Wrapper Tutorial

## Enforce Typesafety at Compile-Time

Sockets cannot detect certain errors at compile-time, e.g.,

```
int s_sd = socket (PF_INET, SOCK_STREAM, 0);
// ...
bind (s_sd, ...); // Bind address.
listen (s_sd); // Make a passive-mode socket.
// Error not detected until run-time.
read (s_sd, buf, sizeof buf);
```

ACE enforces type-safety at compile-time via *factories*, e.g.:

```
SOCK_Acceptor acceptor (port);
// Error: recv() not a method of SOCK_Acceptor.
acceptor.recv (buf, sizeof buf);
```

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## Allow Controlled Violations of Typesafety

*Make it easy to use the C++ Socket wrappers correctly, hard to use it incorrectly, but not impossible to use it in ways the class designers did not anticipate*

- e.g., it may be necessary to retrieve the underlying socket handle:

```
fd_set rd_sds;
FD_ZERO (&rd_sds);
FD_SET (acceptor.get_handle (), &rd_sds);
select (acceptor.get_handle () + 1, &rd_sds, 0, 0, 0);
```

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## Supply Default Parameters

```
SOCK_Connector (SOCK_Stream &new_stream,
                const Addr &remote_sap,
                ACE_Time_Value *timeout = 0,
                const Addr &local_sap = Addr::sap_any,
                int protocol_family = PF_INET,
                int protocol = 0);
```

The result is extremely concise for the common case:

```
SOCK_Stream stream;
// Compiler supplies default values.
SOCK_Connector con (stream, INET_Addr (port, host));
```

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## Define Parsimonious Interfaces

e.g., use LSOCK to pass socket handles:

```
LSOCK_Stream stream;
LSOCK_Acceptor acceptor ("/tmp/foo");
acceptor.accept (stream);
stream.send_handle (stream.get_handle ());
```

versus

```
LSOCK::send_handle (const HANDLE sd) const {
    u_char a[2]; iovec iov; msghdr send_msg;
    a[0] = 0xab, a[1] = 0xcd;
    iov.iov_base = (char *) a; iov.iov_len = sizeof a;
    send_msg.msg_iov = &iov; send_msg.msg_iovlen = 1;
    send_msg.msg_name = (char *) 0;
    send_msg.msg_namelen = 0;
    send_msg.msg_accrighits = (char *) &sd;
    send_msg.msg_accrighitslen = sizeof sd;
    return sendmsg (this->get_handle (), &send_msg, 0);
```

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## Define Parsimonious Interfaces

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```
LSOCK_Stream stream;
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acceptor.accept (stream);
stream.send_handle (stream.get_handle ());
```

versus

```
LSOCK::send_handle (const HANDLE sd) const {
    u_char a[2]; iovec iov; msghdr send_msg;
    a[0] = 0xab, a[1] = 0xcd;
    iov.iov_base = (char *) a; iov.iov_len = sizeof a;
    send_msg.msg_iov = &iov; send_msg.msg_iovlen = 1;
    send_msg.msg_name = (char *) 0;
    send_msg.msg_namelen = 0;
    send_msg.msg_accrighits = (char *) &sd;
    send_msg.msg_accrighitslen = sizeof sd;
    return sendmsg (this->get_handle (), &send_msg, 0);
```

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## Combine Multiple Operations into One Operation

Creating a conventional passive-mode socket requires multiple calls:

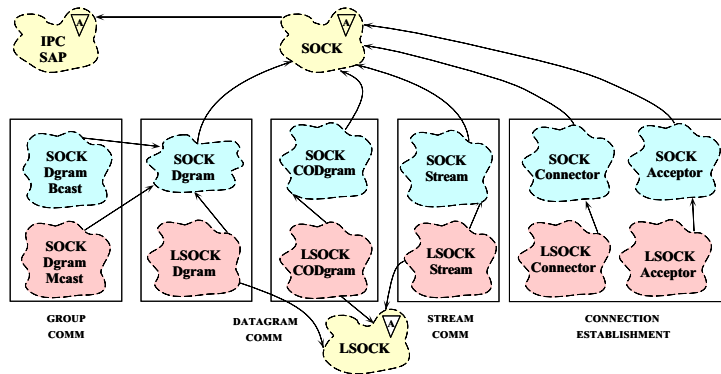
```
int s_sd = socket (PF_INET, SOCK_STREAM, 0);
sockaddr_in addr;
memset (&addr, 0, sizeof addr);
addr.sin_family = AF_INET;
addr.sin_port = htons (port);
addr.sin_addr.s_addr = INADDR_ANY;
bind (s_sd, &addr, addr_len);
listen (s_sd);
// ...
```

SOCK\_Acceptor combines this into a single operation:

```
SOCK_Acceptor acceptor ((INET_Addr) port);
```

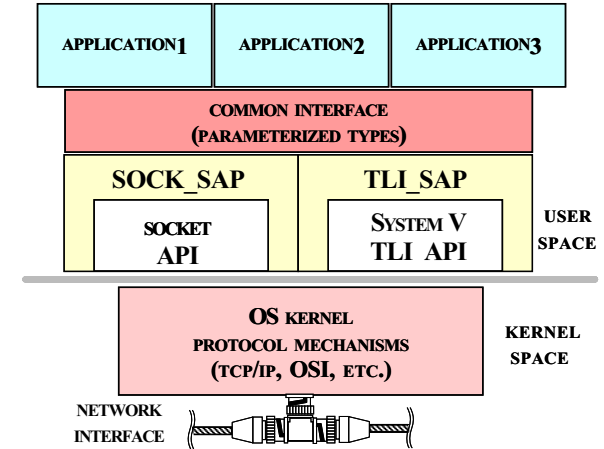
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## Create Hierarchical Class Categories



- Shared behavior is isolated in base classes
- Derived classes implement different communication services, communication domains, and connection roles

## Enhance Portability with Parameterized Types



## Enhance Portability with Parameterized Types (cont'd)

Switching wholesale between sockets and TLI simply requires instantiating a different C++ wrapper, e.g.,

```
// Conditionally select IPC mechanism.
#if defined (USE_SOCKETS)
typedef SOCK_Acceptor PEER_ACCEPTOR;
#elif defined (USE_TLI)
typedef TLI_Acceptor PEER_ACCEPTOR;
#endif // USE_SOCKETS.

int main (void)
{
    // ...

    // Invoke the echo_server with appropriate
    // network programming interfaces.
    echo_server<PEER_ACCEPTOR> (port);
}
```

## Inline Performance Critical Methods

Inlining is time and space efficient since key methods are very short:

```
class SOCK_Stream : public SOCK
{
public:
    ssize_t send (const void *buf, size_t n)
    {
        return ACE_OS::send (this->get_handle (), buf, n);
    }

    ssize_t recv (void *buf, size_t n)
    {
        return ACE_OS::recv (this->get_handle (), buf, n);
    }
};
```

## Define Auxiliary Classes to Hide Error-Prone Details

Standard C socket addressing is awkward and error-prone

- e.g., easy to neglect to zero-out a `sockaddr_in` or convert port numbers to network byte-order, etc.

ACE C++ Socket Wrappers define classes to handle these details

```
class INET_Addr : public Addr {
public:
    INET_Addr (u_short port, long ip_addr = 0) {
        memset (&this->inet_addr_, 0, sizeof this->inet_addr_);
        this->inet_addr_.sin_family = AF_INET;
        this->inet_addr_.sin_port = htons (port);
        memcpy (&this->inet_addr_.sin_addr, &ip_addr, sizeof ip_addr);
    }
    // ...
private:
    sockaddr_in inet_addr_;
};
```

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## Summary of ACE C++ Socket Wrapper Design Principles

- *Domain analysis* identifies and groups related classes of existing API behavior
  - Example *subdomains* include
    - \* Local context management and options, data transfer, connection/termination handling, etc.
    - \* Datagrams vs. streams
    - \* Local vs. remote addressing
    - \* Active vs. passive connection roles
- These relationships are directly reflected in the ACE C++ Socket wrapper inheritance hierarchy

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## Summary of ACE C++ Socket Wrapper Design Principles (cont'd)

- Performance improvements techniques include:
  - Inline functions are used to avoid additional function call penalties
  - Dynamic binding is used sparingly to reduce time/space overhead
    - \* *i.e.*, it is eliminated for `recv/send` path
- Note the difference between the *composition* vs. *decomposition/composition* aspects in design complexity
  - *i.e.*, ACE C++ Socket wrappers are primarily an exercise in *composition* since the basic components already exist
  - More complex OO designs involve both aspects...
    - \* *e.g.*, the ACE Streams, Service Configurator, and Reactor frameworks, etc.

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## Summary of ACE C++ Socket Wrapper Design Principles (cont'd)

- The ACE C++ Socket wrappers are designed to maximize reusability and sharing of components
  - Inheritance is used to *factor out* commonality and *decouple* variation *e.g.*,
    - \* Push common services “upwards” in the inheritance hierarchy
    - \* Factor out variations in client/server portions of socket API
    - \* Decouple datagram vs. stream operations, local vs. remote, etc.
  - Inheritance also supports “functional subsetting”
    - \* *e.g.*, passing open file handles...

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## Concluding Remarks

- Defining C++ wrappers for native OS APIs simplifies the development of correct, portable, and extensible applications
  - C++ `inline` functions ensure that performance isn't sacrificed
- ACE contains many C++ wrappers that encapsulate UNIX, Win32, and RTOS APIs interfaces
  - e.g., sockets, TLI, named pipes, STREAM pipes, etc.
- ACE can be integrated conveniently with CORBA and DCOM provide a flexible high-performance, real-time development framework